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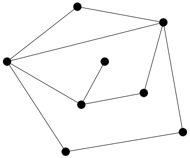
Example. K_4 is planar because there exists a plane drawing of K_4.

Vertices, Edges, and Faces

Definition. In a plane drawing, edges divide the plane into **regions**, or **faces**.



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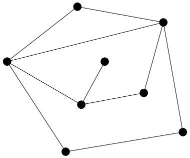

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Notation. Let $p = \#$ of vertices, $q = \#$ of edges, $r = \#$ of regions. Compute the following data:

Graph	p	q	r	
Tetrahedron				
Cube				
Octahedron				
Dodecahedron				
Icosahedron				

In 1750, Euler noticed that ____________ in each of these examples.

Euler's Formula

Theorem 8.1.1 (Euler's Formula) If G is a connected planar graph, then in a plane drawing of G, $p - q + r = 2$.

Proof. (by induction on the number of cycles)

Base Case: If G is a connected graph with no cycles, then G _____

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Let C be a cycle in G, and e be an edge of C. We know that e is adjacent to two different regions, one inside C and one outside C.

Now remove e: Define $H = G \setminus e$. Now H has fewer cycles than G, and one fewer region. The inductive hypothesis holds for H, giving:

Maximal Planar Graphs

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What do we notice about these graphs?

Numerical Conditions on Planar Graphs

- Every face of a maximal planar graph is a triangle!

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Do we need $p \geq 3$?

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Corollary 8.1.3. Every planar graph with $p \geq 3$ vertices has at most $3p - 6$ edges.

- Start with any planar graph G with p vertices and q edges.
- Add edges to G until it is maximal planar. (with $Q \geq q$ edges.)
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Proof.

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Theorem 8.1.7. Every planar graph has a vertex with degree ≤ 5.

Proof.

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Recall: The **girth** $g(G)$ of a graph G is the smallest cycle size.

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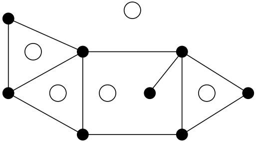
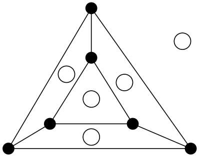
Theorem 8.1.5. If G is planar and bipartite, then $q \leq 2p - 4$.

Theorem 8.1.6. $K_{3,3}$ is not planar.

Dual Graphs

Definition. Given a plane drawing of a planar graph G, the **dual graph** $D(G)$ of G is a graph with vertices corresponding to the regions of G. Two vertices in $D(G)$ are connected by an edge each time the two regions share an edge as a border.





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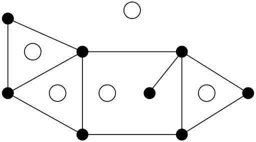

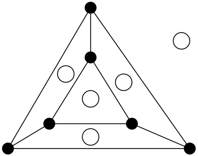

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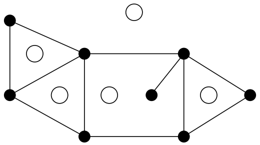

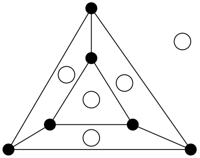

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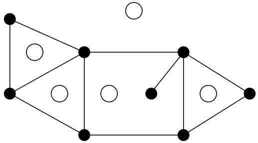

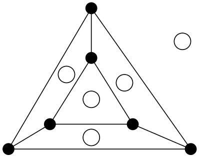

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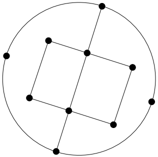
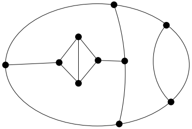
Definition. A graph G is **self-dual** if G is isomorphic to $D(G)$.

Maps

Definition. A *map* is a plane drawing of a connected, bridgeless, planar multigraph. If the map is 3-regular, then it is a **normal map**.





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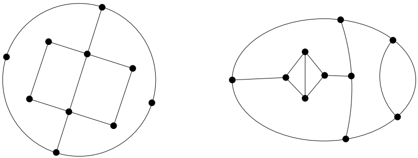


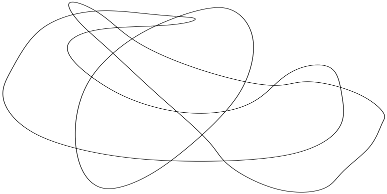
Definition. In a map, the regions are called **countries**. Countries may share several edges.

Definition. A **proper coloring** of a map is an assignment of colors to each country so that no two adjacent countries are the same color.

Question. How many colors are necessary to properly color a map?

Proper Map Colorings

Lemma 8.2.2. If M is a map that is a union of simple closed curves, the regions can be colored by two colors.



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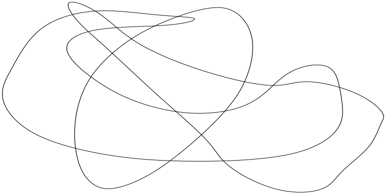

Proof. Color the regions R of M as follows:

$$\begin{Bmatrix} \text{orange} & \text{if } R \text{ is enclosed in an odd number of curves} \\ \text{blue} & \text{if } R \text{ is enclosed in an even number of curves} \end{Bmatrix}.$$

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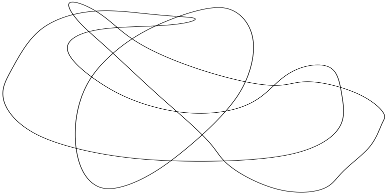

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This is a proper coloring of M. Any two adjacent regions are on opposite sides of a closed curve, so the number of curves in which each is enclosed is off by one.

The Four Color Theorem

Lemma 8.2.6. (The Four Color Theorem)
Every normal map has a proper coloring by four colors.

Proof. Very hard.

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Theorem. If G is a plane drawing of a maximal planar graph, then its dual graph $D(G)$ is a normal map.

- Every face of G is a triangle $\rightsquigarrow$
- G is connected $\rightsquigarrow$
- G is planar $\rightsquigarrow$

The Four Color Theorem

Assuming Lemma 8.2.6,

$$\begin{aligned} G \text{ is maximal planar} &\Rightarrow D(G) \text{ is a normal map} \\ &\Rightarrow \text{countries of } D(G) \text{ 4-colorable} \\ &\Rightarrow \text{vertices of } G \text{ 4-colorable} \\ &\Rightarrow \chi(G) \le 4 \end{aligned}$$

This proves:

Theorem 8.2.8. If G is maximal planar, then $\chi(G) \le 4$.

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Since every planar graph is a subgraph of a maximal planar graph, Lemma C implies:

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