## Edge Coloring

We can also color the edges of a graph.
Definition. An edge coloring of a graph $G$ is a labeling of the edges of $G$ with colors. [Technically, a function $f: E(G) \rightarrow\{1,2, \ldots, /\}$.]

Definition. A proper edge coloring of $G$ is an edge coloring of $G$ such that no two adjacent edges are colored the same.

Example. Cube graph ( $\square_{3}$ ):


We can properly edge color $\square_{3}$ with $\qquad$ colors and no fewer.

Definition. The minimum number of colors necessary to properly edge color a graph $G$ is called the edge chromatic number of $G$, denoted $\chi^{\prime}(G)=$ "chi prime".

## Edge coloring theorems

Question. What is a natural lower bound for $\chi^{\prime}(G)$ ?
Thm 2.2.1: For any graph $G, \chi^{\prime}(G) \geq$ $\qquad$ .
Thm 2.2.2: Vizing's Theorem:
For every graph $G, \chi^{\prime}(G)$ equals either $\Delta(G)$ or $\Delta(G)+1$.
Proof. Hard. (See reference [24] if interested.)
Consequence: To determine $\chi^{\prime}(G)$,

Fact: Most 3-regular graphs have edge chromatic number 3 .


## Snarks

Definition. Another name for 3-regular is cubic.
Definition. A snark is a *bridgeless* cubic graph with edge chromatic number 4.

Example. The Petersen graph $P$ is a snark. It is 3-regular. $\checkmark$
Let us prove that it can not be colored with three colors.
Assume you can color it with three colors. WLOG, assume $a b, a c, a d$.
Either Case 1: be and bi
or Case 2: be and bi.
Either Case 1a: ig and ij or Case 1b: ig and ij.


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In all cases, it is not possible to edge color with 3 colors, so $\chi^{\prime}(G)=4$.

## The edge chromatic number of complete graphs

Goal: Determine $\chi^{\prime}\left(K_{n}\right)$ for all $n$.
Vertex Degree Analysis: The degree of every vertex in $K_{n}$ is $\qquad$ .

Vizing's theorem implies that $\chi^{\prime}\left(K_{n}\right)=$ $\qquad$ or $\qquad$ .

If $\chi^{\prime}\left(K_{n}\right)=\ldots \quad$, then each vertex has an edge leaving of each color.
Question. How many red edges are there?
This is only an integer when:
So, the best we can expect is that $\left\{\begin{array}{l}\chi^{\prime}\left(K_{2 n}\right)= \\ \chi^{\prime}\left(K_{2 n-1}\right)=\end{array}\right.$

The edge chromatic number of complete graphs
Thm 2.2.3: $\quad \chi^{\prime}\left(K_{2 n}\right)=2 n-1 . \quad$ Proof. Use the "turning trick".
Label the vertices of $K_{2 n}$
$0,1, \ldots, 2 n-2, x$.
Connect 0 with $x$,
Connect 1 with $2 n-2$,

Connect $n-1$ with $n$.
Now turn the inside edges.
And do it again. (and again, ...)
Claim: Each turn, new edges are used.


Proof: Each of the edges is a different "circular length".
Vertices are at circular distance $1,3,5, \ldots, 4,2$ from each other, and $x$ is connected to a different vertex each time.

## The edge chromatic number of complete graphs

Theorem 2.2.4: $\quad \chi^{\prime}\left(K_{2 n-1}\right)=2 n-1$.
This construction also gives a way to edge color $K_{2 n-1}$ with $2 n-1$ colors-simply delete vertex $x$ !

This is related to the mathematics of combinatorial designs.
Question. Is it possible for six tennis players to play one match per day in a five-day tournament in such a way that each player plays each other player once?

| Day 1 | $0 x$ | 14 | 23 |
| :--- | :--- | :--- | :--- |
| Day 2 | $1 x$ | 20 | 34 |
| Day 3 | $2 x$ | 31 | 40 |
| Day 4 | $3 x$ | 42 | 01 |
| Day 5 | $4 x$ | 03 | 12 |



Theorem 2.2.3 proves there is such a tournament for all even numbers.

